Segment Routing P4 PoC



Building a PoC of Segment Routing at 100G Using FPGA Smart NIC and P4 Language



Authors

Petr Kaštovský

CEO, Netcope Technologies

Chuck Tato

Director, Communications Business Division Intel® Corporation

Introduction

Service Function Chaining (SFC) is a process of passing network traffic among individual typically virtualized network functions in network function virtualization (NFV) and software-defined network (SDN) infrastructures. The typical SFC can be a series of IDS/IPS, firewalls, wide area network (WAN) optimizers, and load balancers that the traffic needs to go through on its way from client to server and vice versa [1],[2].

One approach to support SFC in NFV infrastructure is to use segment routing, in particular, the IPv6-based segment routing (SRv6) [3]. This technology is now becoming very attractive and deployed in large networks of the future as demonstrated by SoftBank's recent joint announcement with Cisco [4] as well as other deployments also announced by Cisco [5].

During the P4.org workshop in May 2017, teams from Bell Canada, Cisco Systems, and Barefoot Networks presented the concept of The Extensible Network - Evolution in Protocol and Data Plane Agility and explained the benefits of SRv6 for SFC [6].

In our proof of concept, we focused on demonstrating the ability to quickly develop SRv6 acceleration using an Intel® FPGA-based hardware accelerator and P4 programming language [7]. Similarly to accelerating SRv6, other applications can be accelerated using this approach. Good candidates being processing nodes of Vector Packet Processing (VPP) [8].

The theory of segment routing is explained in Figure 1.

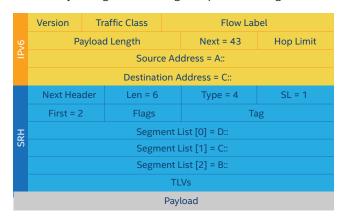


Figure 1. IPv6 and Segment Routing Header. Source [6].

To perform segment routing, SRv6 router goes through the Segment List in Segment Routing Header (SRH) and uses field Segments Left (field SL=1 in Figure 1) as an index of the active segment that is copied over to Destination Address of IPv6 header. We decided to accelerate this data plane operations to demonstrate the productivity and flexibility of P4 language combined with FPGAs.

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P4 Language - Increased Productivity and Abstraction

P4 is a protocol and target independent language that allows for field upgradability of network devices. Furthermore, it significantly improves productivity compared to standard hardware description languages (HDL) that are used to write code for FPGAs. Unlike HDL languages, P4 is domain specific and focused on networking and that makes it an optimal platform for the acceleration of virtual network functions.

Implementation of the SRv6 PoC is quite straightforward. Example 1 shows the P4 code describing packet headers and protocol parser.

```
headers.p4
                                parser.p4
header_type ethernet_t {
                                // General constants
                                #define IPV6_EXT_DEPTH
  fields {
    dstAddr
              : 48:
    srcAddr : 48;
                                // Protocol numbers
                                #define PROTOCOL_IPV6
    etherType: 16;
                                0x86dd
}
                                #define PROTOCOL V6EXT
                                0x2B
header_type ipv6_t {
  fields {
                                // Instances of headers
            : 4;
                                // Outer header stack
    ver
                                metadata seg_meta_t
                                                        lastSeg;
    trafClass: 8;
    flowLab
              : 20;
                                header ethernet_t
                                                     ethernet_0;
    payLen
              : 16;
                                header ipv6_t
    nextHead: 8;
                                ipv6;
    hopLim
                                header ipv6_ext_t
                                                     ipv6_ext;
              : 8;
    srcAddr
              : 128;
                                header ipv6_seg_t
                                                       ipv6_seg;
    dstAddr : 128;
                                // Parse graph
}
                                // Start
                                parser start {
//IPv6, extension header and
                                  return parse_ethernet;
segments
header_type ipv6_ext_t {
  fields {
                                // Parse graph - Outer layers
    nextHead: 8;
                                // Outer ethernet_0
    pad0
           : 16;
                                parser parse_ethernet {
    next_seg : 8;
                                  extract(ethernet_0);
    pad1
            : 32;
                                  return select(latest.etherType) {
                                    PROTOCOL_IPV6
  length: next_seg * 16;
                                parse_ipv6;
}
                                    default
                                : ingress;
header_type ipv6_seg_t {
                                  }
 fields {
    val: 128;
 }
                                parser parse_ipv6 {
}
                                  extract(ipv6);
                                  return select(latest.nextHead)
// Metadata
header_type seg_meta_t {
                                    PROTOCOL_V6EXT: parse_
  fields {
                                ext;
    segVal: 128;
                                    default
                                               : ingress;
    nextSeg: 8;
}
                                parser parse_ext {
                                  extract(ipv6_ext);
                                  return parse_seg;
                                parser parse_seg {
                                  extract(ipv6_seg);
                                  set metadata(lastSeg.
                                segVal,latest.val);
                                  return ingress:
```

Once the packets are parsed and appropriate packet header fields are extracted, it is important to define Match + Action tables that will perform the rewrite of destination IPv6 address by a Segment List item indexed by the Segments Left field and decrement of the Segments Left field.

```
tables.p4
                                  main.p4
// Actions
                                   #include "headers.p4"
action rewrite() {
                                   #include "parser.p4"
                                   #include "tables.p4"
  modify_field(ipv6.
dstAddr,lastSeg.segVal);
                                   control ingress {
  add_to_field(ipv6_ext.
                                    if(valid(ipv6 ext)) {
next_seg, -1);
                                       apply(tab_rewrite);
                                  }
// Tables
// Table that does nothing
table tab_rewrite {
  actions {
    rewrite:
  }
}
```

Example 2. SRv6 Processing Code in P4

FPGA-Based Smart NIC as a Target

P4, being a target independent language, makes it easy for the designer to select Intel FPGAs as the target architecture. Intel FPGAs are a preferred hardware platform option for its efficiency when compared to CPUs and NPUs.

Using a P4 language as a programming language waives that barriers that are traditionally associated with programming FPGAs and brings the benefits of FPGAs into the networking domain.

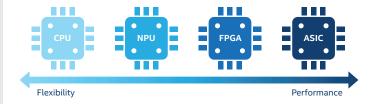


Figure 2. Netcope P4 Improves FPGA Flexibility

Netcope projects:

 Netcope P4 - Netcope P4 is an FPGA-vendor independent project by Netcope Technologies providing integration into different flavors of FPGA-based smart network interface cards (NICs) offering up to 2X 100GbE network capacity to fully deliver on the improved efficiency over NPUs.[†]

https://www.netcope.com/en/products/netcopep4

Executing the PoC

To perform the PoC we have used the online Netcope P4 Cloud to perform the synthesis of P4 code. The whole compilation process takes several steps but all of them are simplified for the user to only upload P4 code at the beginning and get FPGA bitstream at the end.

Step 1:

After uploading the code, an email confirmation will be sent confirming the task being queued for compilation and synthesis. The time required for compilation and synthesis is around two and a half hours for this PoC. Another email will be sent confirming that the task has finished. Also, it is possible to check the status of the task on the portal shown in Figure 3.

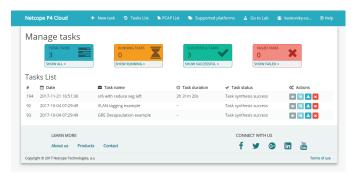


Figure 3. Netcope P4 Cloud GUI

Step 2:

Test the generated FPGA bitstream in the real hardware with the Ethernet cable connected in loopback. All that is needed is to load the bitstream into the FPGA, configure P4 Match&Action tables, start packet capture, and send test traffic from packet capture (PCAP) files. Figure 4 is an example of the specific packet before and after PoC SRv6 pipeline, which was dissected by Wireshark.

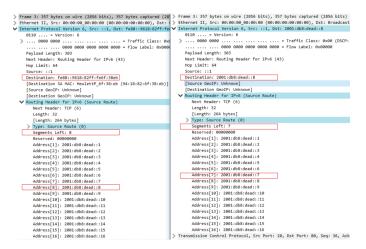


Figure 4. Packet Content Before and After SRv6 Processing

The code is functionally correct. Using the Intel Arria® 10 FPGA on the pre-production version of the Intel FPGA Programmable Acceleration Card N3000 with 8x10GE interfaces, we were able to achieve from 78 to 80 Gbps of port to port throughput without any specific performance optimizations or tunings. See Figure 5.

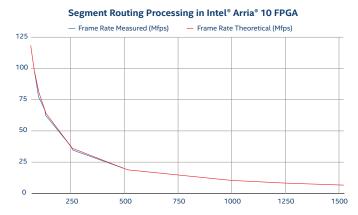


Figure 5. Throughput Achieved with the Real Hardware

It was a very interesting exercise where most of the time to get the PoC working was figuring out the problem on the algorithmic side. In other words, what data transformation should be done. Once this was clear, writing the code was a matter of day and getting the bitstream including test in hardware was a matter of 3 hours. This is a true revolution for Intel FPGA-based NIC programming.

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[†] Tests measure performance of components on a particular test, in specific systems. Differences in hardware, software, or configuration will affect actual performance. Consult other sources of information to evaluate performance as you consider your purchase. For more complete information about performance and benchmark results, visit www.intel.com/benchmarks.

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